

Introduction to Quantum Algorithms-Lecture 4

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Problem

From a list of $q = 2^n$ items to peak the k -th item.

Example from a phone book with 10^6 names to find the k -th name.

Mathematical Introduction

Let $|\sigma\rangle$ and $|s\rangle$ two perpendicular unity vectors

$$\langle \sigma | s \rangle = \langle s | \sigma \rangle = 0, \quad \langle \sigma | \sigma \rangle = \langle s | s \rangle = 1, \quad |\tau\rangle = \sin \omega |\sigma\rangle + \cos \omega |s\rangle$$

Matrix representation

$$|\sigma\rangle \leftrightarrow \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \quad |s\rangle \leftrightarrow \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \quad |\tau\rangle \leftrightarrow \begin{pmatrix} \sin \omega \\ \cos \omega \end{pmatrix},$$

Grover or rotation operator

$$G = (2|\tau\rangle\langle\tau| - \mathbb{I})(\mathbb{I} - 2|\sigma\rangle\langle\sigma|) \leftrightarrow \begin{pmatrix} \cos 2\omega & \sin 2\omega \\ -\sin 2\omega & \cos 2\omega \end{pmatrix},$$

$$G^m \leftrightarrow \begin{pmatrix} \cos 2m\omega & \sin 2m\omega \\ -\sin 2m\omega & \cos 2m\omega \end{pmatrix},$$

$$|\rho\rangle = \sin \theta |\sigma\rangle + \cos \theta |s\rangle \rightsquigarrow G^m |\rho\rangle \leftrightarrow \begin{pmatrix} \cos(2\omega + \theta) \\ \sin(2\omega + \theta) \end{pmatrix}$$

$$|\sigma\rangle = |\bar{k}\rangle, |\tau\rangle = \frac{1}{\sqrt{q}} \sum_{\ell=0}^{q-1} |\bar{\ell}\rangle \rightsquigarrow |s\rangle = \frac{1}{\sqrt{q-1}} \sum_{\ell=0, \ell \neq k}^{q-1} |\bar{\ell}\rangle$$

$$|\tau\rangle = \frac{1}{\sqrt{q}} |\sigma\rangle + \sqrt{\frac{q-1}{q}} |s\rangle \rightsquigarrow \sin \omega = \frac{1}{\sqrt{q}}$$

$$|\rho\rangle = |\tau\rangle = \sin \omega |\sigma\rangle + \cos \omega |s\rangle$$

$$G^m |\rho\rangle = \sin((2m+1)\omega) |\sigma\rangle + \cos((2m+1)\omega) |s\rangle$$

Approximation

If $(2m+1)\omega \sim \frac{\pi}{2}$ then $G^m |\rho\rangle \sim |\sigma\rangle$

$$\sin \omega = \frac{1}{\sqrt{q}} \rightsquigarrow \omega \sim \frac{1}{\sqrt{q}} \Rightarrow (2m+1) \frac{1}{\sqrt{q}} \sim \frac{\pi}{2} \rightsquigarrow m \sim \frac{\pi}{2} \sqrt{q}$$

n	$q = 2^n$	m	Prob
10	1024	24	99.85%
20	10^6	803	99.999%

Grover's Algorithm

Step 1 : We start with the initial state $|\text{in}\rangle = |0\rangle = |00\dots 0\rangle$

Step 2 : We apply the operator $H^{\otimes n}$ on the initial state and we get the state $|\rho\rangle = \frac{1}{\sqrt{q}} \sum_{\ell=0}^{q-1} \left| \ell \right\rangle$

Step 3 : We apply $m = \left[\left(\frac{\pi}{2\omega} - 1 \right) / 2 \right]$ times the Grover operator, with $\omega = \arcsin \frac{1}{\sqrt{q}}$

$$G = (2|\tau\rangle\langle\tau| - \mathbb{I})(\mathbb{I} - 2|\sigma\rangle\langle\sigma|) \leftrightarrow \begin{pmatrix} \cos 2\omega & \sin 2\omega \\ -\sin 2\omega & \cos 2\omega \end{pmatrix},$$

The final state is given by the formula

$$|\text{fin}\rangle G^m |\rho\rangle = \sin((2m+1)\omega) |\sigma\rangle + \cos((2m+1)\omega) |s\rangle$$

Step 4 : The probability to be in the final state $|\sigma\rangle$ is equal to $\sin^2((2m+1)\omega)$